## Oblivious Transfer

## CS 598 DH

## Today's objectives

Review semi-honest security Introduce oblivious transfer (OT)

Build OT from DDH

See an end-to-end security proof


## Two-Party Semi-Honest Security for deterministic functionalities

Let $f_{0}, f_{1}$ be functions. We say that a protocol $\Pi$ securely computes $f_{0}, f_{1}$ in the presence of a semi-honest adversary if for each party $i \in\{0,1\}$ there exists a polynomial time simulator $\mathcal{S}_{i}$ such that for all inputs $x_{0}, x_{1}$ :
$\operatorname{View}_{i}^{\Pi}\left(x_{0}, x_{1}\right) \stackrel{c}{=} \mathcal{S}_{i}\left(x_{i}, f_{i}\left(x_{0}, x_{1}\right)\right)$

## Semi-honest Security



Three notions of "hard to tell apart"
$X \equiv Y \quad$ Identically distributed
$X \approx Y \quad$ Statistically close
As we increase a parameter, the distributions quickly become close together.
$X \stackrel{c}{=} Y \quad$ Indistinguishable As we increase a parameter, it quickly becomes difficult for programs to tell the distributions apart.

## Oblivious Transfer



1-out-of-2
Oblivious Transfer


Receiver


Receiver



## 1-out-of-2 OT Ideal Functionality

## $m_{0}, m_{1}$ <br> $$
b \in\{0,1\}
$$



Receiver


OT is an extremely powerful tool
Given enough OTs, we can build a semi-honest protocol for any computable function


## Secure AND

## $0, x \quad y$ <br> OT

## Secure AND

$$
\begin{aligned}
& \xrightarrow{0, x} \underset{ }{\longleftrightarrow} \quad y \\
& \text { OT } \\
& y \\
& \left(\left\{\begin{array}{ll}
0 & \text { if } y=0 \\
x & \text { if } y=1
\end{array}\right)=x \wedge y\right.
\end{aligned}
$$



## Public Key Encryption Scheme

Generating a key makes a public key, private key pair pk, sk
Anyone with pk can encrypt messages
Only those with sk can decrypt

## Intuitive Idea for OT

Receiver makes two public keys, but only one has a matching private key

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Receiver makes two public keys, but only one has a matching private key Receiver sends each public key to Sender

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Sender encrypts one message per key

## Intuitive Idea for OT

Receiver makes two public keys, but only one has a matching private key Receiver sends each public key to Sender

Sender encrypts one message per key
Receiver decrypts (only) the desired message

Goal:

## Correctness

Semi-honest Security


Goal:

## Correctness

Semi-honest Security

$$
\begin{gathered}
\operatorname{View}_{S}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right) \approx \mathcal{S}_{S}\left(m_{0}, m_{1}, \perp\right) \\
\operatorname{View}_{R}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right) \approx \mathcal{S}_{R}\left(b, m_{b}\right)
\end{gathered}
$$

## Decisional Diffie-Hellman Assumption

"It is hard to compute logarithms in certain mathematical sets"

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"It is hard to compute logarithms in certain mathematical sets"

## Let $G$ be a cyclic group of order $q$ with generator $g$



$$
\begin{aligned}
& \text { Ideal }(\quad): \\
& \quad a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad b \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad c \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{g^{a}, g^{b}, g^{c}\right\}
\end{aligned}
$$

$m_{0}, m_{1}$

Sender

$$
\begin{aligned}
& a \stackrel{S}{\leftarrow} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \stackrel{\$}{\leftarrow} G
\end{aligned}
$$

Receiver
$m_{0}, m_{1}$

Sender

$$
\begin{aligned}
& a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a}
\end{aligned}
$$

$$
h_{0}, h_{1}
$$

$m_{0}, m_{1}$

Sender

$$
h_{0}, h_{1}
$$

$$
\begin{aligned}
& a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \stackrel{\$}{\leftarrow} G
\end{aligned}
$$

Receiver

$$
\begin{aligned}
& r_{0} \stackrel{\$}{\stackrel{\$}{\leftarrow} \mathbb{Z}_{q}} \\
& r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}
\end{aligned}
$$

$m_{0}, m_{1}$

Sender

$$
h_{0}, h_{1}
$$

$$
\begin{aligned}
& a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \stackrel{\$}{\leftarrow} G
\end{aligned}
$$

$$
\begin{aligned}
& r_{0} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}
\end{aligned}
$$


$m_{0}, m_{1}$

Sender

$$
h_{0}, h_{1}
$$

$$
\begin{aligned}
& r_{0} \stackrel{\$}{\stackrel{\$}{\leftarrow}} \mathbb{Z}_{q} \\
& r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}
\end{aligned}
$$


$m_{0}, m_{1}$

$$
a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}
$$

Receiver
Sender $\quad r_{0} \stackrel{\stackrel{\S}{\leftarrow} \mathbb{Z}_{q}}{ } \quad r_{1} \stackrel{\lessgtr}{\leftrightarrows} \mathbb{Z}_{q}$

$\frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}$
$m_{0}, m_{1}$

$$
\begin{aligned}
& a \stackrel{\$}{\gtrless} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b}{ }^{g} G
\end{aligned}
$$

$$
h_{0}, h_{1}
$$

Sender

$$
\begin{aligned}
& r_{0} \stackrel{\stackrel{8}{\leftarrow} \mathbb{Z}_{q}}{\stackrel{\$}{\leftarrow} \mathbb{Z}_{q}}
\end{aligned}
$$

$$
\xrightarrow{\begin{array}{ll}
g^{r_{0}} \\
h_{0}^{r_{0}} \cdot m_{0}
\end{array}} \begin{aligned}
& h_{1}^{r_{1}} \cdot m_{1}
\end{aligned} \longrightarrow \frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}
$$

Receiver

$$
\frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}=\frac{\left(g^{a}\right)^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}
$$

$m_{0}, m_{1}$
$a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$

Receiver
Sender $\quad r_{0} \stackrel{\stackrel{\S}{\leftarrow} \mathbb{Z}_{q}}{ } \quad r_{1} \stackrel{\lessgtr}{\leftrightarrows} \mathbb{Z}_{q}$

$$
\xrightarrow{\begin{array}{ll}
g^{r_{0}} \\
h_{0}^{r_{0}} \cdot m_{0}
\end{array}} \begin{aligned}
& h_{1}^{r_{1}} \cdot m_{1}
\end{aligned} \longrightarrow \frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}
$$

$$
\frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}=\frac{\left(g^{a}\right)^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}=\frac{g^{a \cdot r_{b}} \cdot m_{b}}{g^{a \cdot r_{b}}}
$$

$m_{0}, m_{1}$

$$
\begin{aligned}
& a \stackrel{\$}{\gtrless} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \stackrel{\Phi}{\leftarrow} G
\end{aligned}
$$

$$
h_{0}, h_{1}
$$

Sender

$$
\begin{aligned}
& r_{0} \stackrel{\stackrel{\leftrightarrow}{\leftarrow} \mathbb{Z}_{q}}{\stackrel{\&}{\leftarrow}} \begin{array}{l}
\mathbb{Z}_{q}
\end{array}
\end{aligned}
$$

Receiver

$$
\frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}=\frac{\left(g^{a}\right)^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}=\frac{g^{a \cdot r_{b}} \cdot m_{b}}{g^{a \cdot r_{b}}}=m_{b}
$$

$m_{0}, m_{1}$


$$
a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}
$$

Receiver
Sender $\begin{array}{ll}r_{0} & \stackrel{\S}{\leftarrow} \mathbb{Z}_{q} \\ & r_{1} \\ & \end{array}$

$$
\xrightarrow{\begin{array}{ll}
g^{r_{0}} \\
h_{0}^{r_{0}} \cdot m_{0}
\end{array} h_{1}^{h_{1}^{r_{1}} \cdot m_{1}}} \longrightarrow \frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}
$$

$\operatorname{View}_{S}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right)=\cdots$
$m_{0}, m_{1}$
$a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$

Receiver
Sender $\quad r_{0} \stackrel{\$ \not \mathbb{Z}_{q}}{ } \quad r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$


$$
\begin{gathered}
\operatorname{View}_{S}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right)=\left\{m_{0}, m_{1}, h_{0}, h_{1}, r_{0}, r_{1}\right\} \\
\bar{\equiv} \\
\mathcal{S}_{S}\left(m_{0}, m_{1}, \perp\right): \\
h_{0}, h_{1}, r_{0}, r_{1} \stackrel{\$}{\leftarrow} G \\
\operatorname{return}\left\{m_{0}, m_{1}, h_{0}, h_{1}, r_{0}, r_{1}\right\}
\end{gathered}
$$

$m_{0}, m_{1}$

$$
\begin{aligned}
& r_{0} \stackrel{\&}{\leftarrow} \mathbb{Z}_{q} \\
& r_{1} \stackrel{\&}{\&} \mathbb{Z}_{q}
\end{aligned}
$$

Sender $\quad r_{0} \stackrel{\Phi}{\leftarrow} \mathbb{Z}_{q}$

$$
a \stackrel{\$}{\gtrless} \mathbb{Z}_{q}
$$

$$
h_{0}, h_{1}
$$

$$
h_{b} \leftarrow g^{a}
$$

$$
h_{1-b}^{\stackrel{\circ}{\&}} G
$$

$$
\left.\xrightarrow{\begin{array}{ll}
g^{r_{0}} \\
h_{0}^{r_{0}} \cdot m_{0}
\end{array}} \begin{array}{l}
h_{1}^{r_{1}} \cdot m_{1}
\end{array}\right] \frac{h_{b}^{r_{b}} \cdot m_{b}}{\left(g^{r_{b}}\right)^{a}}
$$

$$
\operatorname{View}_{R}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right)=\left\{b, a, h_{1-b}, g^{r_{0}}, g^{r_{1}}, h_{b}^{r_{b}} \cdot m_{b}, h_{1-b}^{r_{1-b}} \cdot m_{1-b}\right\}
$$

$$
\begin{aligned}
& \mathcal{S}_{R}\left(b, m_{b}\right): \\
& \quad r_{0}, r_{1}, a, k, s \stackrel{\Phi}{\leftarrow} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, g^{a \cdot r_{b}} \cdot m_{b}, g^{s}\right\}
\end{aligned}
$$

$m_{0}, m_{1}$

$$
\begin{aligned}
& a \stackrel{\$}{\gtrless} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \$^{\&} G
\end{aligned}
$$

$$
\begin{array}{lll}
r_{0} \stackrel{\S}{\stackrel{~}{\gtrless}} \mathbb{Z}_{q} \\
r_{1} \stackrel{\&}{\leftarrow} \\
\end{array}
$$

Sender $\quad r_{0} \stackrel{\Phi}{\leftarrow} \mathbb{Z}_{q}$

$$
h_{0}, h_{1}
$$

$$
\operatorname{View}_{R}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right)=\left\{b, a, h_{1-b}, g^{r_{0}}, g^{r_{1}}, h_{b}^{r_{b}} \cdot m_{b}, h_{1-b}^{r_{1-b}} \cdot m_{1-b}\right\}
$$

$$
\begin{aligned}
& \mathcal{S}_{R}\left(b, m_{b}\right): \\
& \quad r_{0}, r_{1}, a, k, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, g^{a \cdot r_{b}} \cdot m_{b}, g^{s}\right\}
\end{aligned}
$$

$\operatorname{View}_{R}^{\mathrm{OT}}\left(m_{0}, m_{1}, b\right)=\left\{b, a, h_{1-b}, g^{r_{0}}, g^{r_{1}}, h_{b}^{r_{b}} \cdot m_{b}, h_{1-b}^{r_{1-b}} \cdot m_{1-b}\right\}$
"DDH implies that $h_{1-b}^{r_{1-b}}$ "looks random", and

$$
h_{1-b}^{r_{1-b}} \text { masks message } m_{1-b} "
$$

$$
\begin{aligned}
& \mathcal{S}_{R}\left(b, m_{b}\right): \\
& \quad r_{0}, r_{1}, a, k, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, g^{a \cdot r_{b}} \cdot m_{b}, g^{s}\right\}
\end{aligned}
$$

$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{b} \leftarrow g^{a} \\
& h_{1-b} \stackrel{\Phi}{\leftarrow} G \\
& \text { return }\left\{b, a, h_{1-b}, g^{r_{0}}, g^{r_{1}}, h_{b}^{r_{b}} \cdot m_{b}, h_{1-b}^{r_{1-b}} \cdot m_{1-b}\right\}
\end{aligned}
$$

$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\S}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\Im}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \$ \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\$}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

R's input
$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \$ \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\$}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

R's input


R's randomness
$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\$}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

R's input


R's randomness
S's random messages
$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\mathbb{Z}}{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\S}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

R's input


R's randomness

S's random messages

The message R can decrypt
$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :
WLOG, suppose $b=0$

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \stackrel{\oiint}{\leftarrow} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{b}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

R's input


R's randomness

S's random messages

The message
R cannot decrypt

The message R can decrypt
$\qquad$
$\operatorname{Hyb0}\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftrightarrows G \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

$$
\operatorname{Hyb1}\left(m_{0}, m_{1}, b\right):
$$

$$
a, r_{0}, r_{1}, k \stackrel{\mathbb{Z}_{q}}{ }
$$

$$
h_{0} \leftarrow g^{a}
$$

$$
h_{1} \leftarrow g^{k}
$$

$$
\text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, h_{1}^{r_{1}} \cdot m_{1}\right\}
$$

$$
\begin{aligned}
& \operatorname{Hyb1}\left(m_{0}, m_{1}, b\right): \\
& \quad a, r_{0}, r_{1}, k \stackrel{\bigotimes}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& \quad \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, h_{1}^{r_{1}} \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb} 2\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k \stackrel{\mathbb{Z}_{q}}{ } \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& \text { mask } \leftarrow h_{1}^{r_{1}} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\} \\
& \quad=
\end{aligned}
$$

$$
\operatorname{Hyb1}\left(m_{0}, m_{1}, b\right):
$$

$$
a, r_{0}, r_{1}, k \stackrel{\mathbb{Z}}{q}
$$

$$
h_{0} \leftarrow g^{a}
$$

$$
h_{1} \leftarrow g^{k}
$$

$$
\text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, h_{1}^{r_{1}} \cdot m_{1}\right\}
$$

$\operatorname{Hyb} 2\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0}, r_{1}, k \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& \text { mask } \leftarrow h_{1}^{r_{1}} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

$$
\begin{aligned}
& \operatorname{Hyb} 2\left(m_{0}, m_{1}, b\right): \\
& \quad a, r_{0}, r_{1}, k \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& \text { mask } \leftarrow h_{1}^{r_{1}} \\
& \quad \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\} \\
& \text { Hyb3 }\left(m_{0}, m_{1}, b\right): \\
& \quad a, r_{0}, r_{1}, k \stackrel{\$}{\&} \mathbb{Z}_{q} \\
& \quad h_{0} \leftarrow g^{a} \\
& \quad h_{1} \leftarrow g^{k} \\
& g^{\prime} \leftarrow g^{r_{1}} \\
& \text { mask } \leftarrow g^{k \cdot r_{1}} \\
& \quad \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g_{51}^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb} 3\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0}, r_{1}, k \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& g^{\prime} \leftarrow g^{r_{1}} \\
& \text { mask } \leftarrow g^{k \cdot r_{1}} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb} 3\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0}, r_{1}, \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& g^{\prime} \leftarrow g^{r_{1}} \\
& \text { mask } \leftarrow g^{k \cdot r_{1}} \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb4}\left(m_{0}, m_{1}, b\right)$ :
$a, r_{0} \stackrel{\S}{\leftarrow} \mathbb{Z}_{q}$
$h_{0} \leftarrow g^{a}$
$\left\{h_{1}, g^{\prime}\right.$, mask $\} \leftarrow \operatorname{Real}()$
return $\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}\right.$, mask $\left.\cdot m_{1}\right\}$

## Decisional Diffie-Hellman Assumption

"It is hard to compute logarithms in certain mathematical sets"

## Let $G$ be a cyclic group of order $q$ with generator $g$



$$
\begin{aligned}
& \text { Ideal }(\quad): \\
& \quad a \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad b \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad c \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{g^{a}, g^{b}, g^{c}\right\}
\end{aligned}
$$

$\operatorname{Hyb4}\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \left\{h_{1}, g^{\prime}, \operatorname{mask}\right\} \leftarrow \operatorname{Real}() \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \operatorname{mask} \cdot m_{1}\right\}
\end{aligned}
$$

Real():
$k, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$
return $\left\{g^{k}, g^{r_{1}}, g^{k \cdot r_{1}}\right\}$
$\operatorname{Hyb4}\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \left\{h_{1}, g^{\prime}, \operatorname{mask}\right\} \leftarrow \operatorname{Real}() \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

## $\underline{C} \quad[B y D D H]$

$\operatorname{Hyb} 5\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \left\{h_{1}, g^{\prime}, \text { mask }\right\} \leftarrow \text { Ideal }() \\
& \text { return }\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}, \text { mask } \cdot m_{1}\right\}
\end{aligned}
$$

## Real():

$k, r_{1} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$
return $\left\{g^{k}, g^{r_{1}}, g^{k \cdot r_{1}}\right\}$

Ideal():
$k, r_{1}, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$
return $\left\{g^{k}, g^{r_{1}}, g^{s}\right\}$
$\operatorname{Hyb5}\left(m_{0}, m_{1}, b\right)$ :
$a, r_{0} \stackrel{\&}{\leftarrow} \mathbb{Z}_{q}$
$h_{0} \leftarrow g^{a}$
$\left\{h_{1}, g^{\prime}\right.$, mask $\} \leftarrow \operatorname{Ideal}()$
return $\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}\right.$, mask $\left.\cdot m_{1}\right\}$

## Ideal():

$k, r_{1}, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$
return $\left\{g^{k}, g^{r_{1}}, g^{s}\right\}$
$\operatorname{Hyb5}\left(m_{0}, m_{1}, b\right)$ :
$a, r_{0} \stackrel{\$}{\leftarrow} \mathbb{Z}_{q}$
$h_{0} \leftarrow g^{a}$
$\left\{h_{1}, g^{\prime}\right.$, mask $\} \leftarrow$ Ideal ( )
return $\left\{b, a, h_{1}, g^{r_{0}}, g^{\prime}, h_{0}^{r_{0}} \cdot m_{0}\right.$, mask $\left.\cdot m_{1}\right\}$

$$
\begin{aligned}
& \text { Ideal( ): } \\
& \quad k, r_{1}, s \stackrel{\&}{\gtrless} \mathbb{Z}_{q} \\
& \quad \text { return }\left\{g^{k}, g^{r_{1}}, g^{s}\right\}
\end{aligned}
$$

$\operatorname{Hyb5}\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{乌}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& g^{\prime} \leftarrow g^{r_{1}} \\
& \text { mask } \leftarrow g^{s}
\end{aligned}
$$

$$
\text { return }\left\{b, a, h_{1}, g^{r_{0}}, g_{s_{0}^{\prime}}^{\prime} h_{0}^{r_{0}} \cdot m_{0} \text {, mask } \cdot m_{1}\right\}
$$

$\operatorname{Hyb5}\left(m_{0}, m_{1}, b\right):$

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \leftleftarrows \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& h_{1} \leftarrow g^{k} \\
& g^{\prime} \leftarrow g^{r_{1}} \\
& \text { mask } \leftarrow g^{s}
\end{aligned}
$$

return $\left\{b, a, h_{1}, g^{r_{0}}, g_{g_{9}^{\prime}}^{\prime} h_{0}^{r_{0}} \cdot m_{0}\right.$, mask $\left.\cdot m_{1}\right\}$

Hyb6( $\left.m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s} \cdot m_{1}\right\}
\end{aligned}
$$

## $=$

$\operatorname{Hyb5}\left(m_{0}, m_{1}, b\right)$ :

$$
a, r_{0}, r_{1}, k, s \stackrel{\mathbb{Z}_{q}}{ }
$$

$$
h_{0} \leftarrow g^{a}
$$

$$
h_{1} \leftarrow g^{k}
$$

$$
g^{\prime} \leftarrow g^{r_{1}}
$$

$$
\text { mask } \leftarrow g^{s}
$$

return $\left\{b, a, h_{1}, g^{r_{0}}, g_{e_{0}^{\prime}}^{\prime} h_{0}^{r_{0}} \cdot m_{0}\right.$, mask $\left.\cdot m_{1}\right\}$

Hyb6( $\left.m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \$ \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s} \cdot m_{1}\right\}
\end{aligned}
$$

$\operatorname{Hyb6}\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{\mathbb{Z}_{q}}{ } \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s} \cdot m_{1}\right\}
\end{aligned}
$$

## = [By one-time-pad]

$\operatorname{Hyb7}\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{\S}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s}\right\}
\end{aligned}
$$

$\operatorname{Hyb} 7\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{\S}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s}\right\}
\end{aligned}
$$

$$
\begin{aligned}
& \mathcal{S}_{R}\left(b, m_{0}\right): \\
& \quad a, r_{0}, r_{1}, k, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& \quad h_{0} \leftarrow g^{a} \\
& \quad \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s}\right\}
\end{aligned}
$$

$$
\overline{=}
$$

Hyb7 $\left(m_{0}, m_{1}, b\right)$ :

$$
\begin{aligned}
& a, r_{0}, r_{1}, k, s \stackrel{\$}{\leftarrow} \mathbb{Z}_{q} \\
& h_{0} \leftarrow g^{a} \\
& \text { return }\left\{b, a, g^{k}, g^{r_{0}}, g^{r_{1}}, h_{0}^{r_{0}} \cdot m_{0}, g^{s}\right\}
\end{aligned}
$$

## Today's objectives

Review semi-honest security Introduce oblivious transfer (OT)

Build OT from DDH

See an end-to-end security proof

